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(Joint work with Ventsislav Chonev and Joël Ouaknine)

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Reachability for Continuous-Time Markov Chains



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Distribution P(t) at time t satisfies P'(t) = P(t)Q, where

$$Q = \begin{pmatrix} -0.025 & 0.02 & 0.005 \\ 0.3 & -0.5 & 0.2 \\ 0.02 & 0.4 & -0.42 \end{pmatrix}$$

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• Reduce to the **time-bounded case** by computing the stationary distribution:

$$\pi = (0.885, 0.071, 0.044)$$

• Require that π not be on boundary of the target set.

"To analyze a cyber-physical system, such as a pacemaker, we need to consider the **discrete software controller** interacting with the physical world, which is typically modelled by **differential equations**"

Rajeev Alur (CACM, 2013)



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Is ever more likely to be a Bear market than a Bull market:

 $\exists t \left(P(t)_{\text{Bear}} \geq P(t)_{\text{Bull}} \right) ?$

$$\begin{aligned} \mathbf{x} &: \mathbb{R}_{\geq 0} \to \mathbb{R}^k \\ \dot{\mathbf{x}} &= \mathbf{A} \mathbf{x} \end{aligned}$$

$$\begin{split} \mathbf{x} &: \mathbb{R}_{\geq 0} \to \mathbb{R}^k \\ \dot{\mathbf{x}} &= \mathbf{A} \mathbf{x} \\ \Rightarrow & \mathbf{x}(t) = \exp(\mathbf{A} t) \mathbf{x}(0) \end{split}$$



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Note – the λ_j are complex in general.

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BOUNDED-ZERO Problem

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• Decidability open! [Bell, Delvenne, Jungers, Blondel 2010]

A lot of work since 1920s on the zeros of exponential polynomials

$$f(z) = \sum_{j=1}^m P_j(z) e^{\lambda_j z}$$

(Polya, Ritt, Tamarkin, Kac, Voorhoeve, van der Poorten, ...) but mostly on distribution of *complex* zeros.

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CONTINUOUS-ORBIT Problem

The problem of whether the trajectory $\mathbf{x}(t) = e^{\mathbf{A}t}\mathbf{x}(0)$ reaches a given target point was shown to be decidable by Hainry (2008) and in PTIME by Chen, Han and Yu (2015).

Theorem (Chonev, Ouaknine, W. 2015)

Assuming Schanuel's Conjecture, BOUNDED-ZERO is decidable at all orders.

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It turns out that decidability in the bounded case follows from a much more general result, discovered (but not published) in the early 1990s by Macintyre and Wilkie.

[Angus Macintyre, personal communication, July 2015]

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Schanuel's Conjecture

If z_1, \ldots, z_n are complex numbers linearly independent over \mathbb{Q} then some *n*-element subset of $\{z_1, \ldots, z_n, e^{z_1}, \ldots, e^{z_n}\}$ is algebraically independent.



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Easy Consequence

By Schanuel's conjecture, some two-element subset of $\{1, \pi i, e^1, e^{\pi i}\}$ is algebraically independent.

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Theorem (Macintyre and Wilkie 1996)

The first-order theory of $(\mathbb{R}, +, \cdot, e^x)$ is decidable, assuming Schanuel's conjecture.









'non-trivial' zero \Rightarrow t^{*} transcendental



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Real-valued exponential polynomial $f(t) = \sum_{j=1}^{m} P_j(t) e^{\lambda_j t}$



Can this situation arise?



Easily! For example, $f(t) = 2 + e^{it} + e^{-it}$.

Example

• Write $f(t) = 2 + e^{it} + e^{-it}$ in the form $f(t) = P(e^{it})$ for the Laurent polynomial

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• Factorisation $P(z) = (1 + z)(1 + z^{-1})$ induces a factorisation

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Idea: factorise f. Noting that factors may be complex-valued!

Any exponential polynomial f(t) can be written

$$f(t) = P(t, e^{a_1 t}, \ldots, e^{a_m t})$$

with

$$P \in \mathbb{C}[x, x_1^{\pm 1}, \dots, x_m^{\pm 1}]$$

and $\{a_1, \ldots, a_m\}$ a set of complex algebraic numbers linearly independent over \mathbb{Q} .

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Proof Strategy

Factorisation of P into irreducible factors induces factorisation of f. Assuming Schanuel's conjecture, we can decide the existence of zeros of real-valued and complex-valued irreducible factors.

ZERO Problem

Instance: fQuestion: Is there $t \in \mathbb{R}_{\geq 0}$ such that f(t) = 0?

How well can one approximate a real number x with rationals?

$$\left|x-\frac{p}{q}\right|$$

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Theorem (Roth 1955)

Let $x \in \mathbb{R}$ be algebraic. Then for any $\varepsilon > 0$ there are finitely many integers p, q such that

$$\left|x-\frac{p}{q}\right|<rac{1}{q^{2+arepsilon}}\,.$$



How well can one approximate a real number x with rationals?

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Definition

Let $x \in \mathbb{R}$. The **Diophantine-approximation type** L(x) is:

$$L(x) = \inf \left\{ c : \left| x - rac{p}{q} \right| < rac{c}{q^2} ext{ for some } p, q \in \mathbb{Z}
ight\}.$$

Continued Fractions

Finite continued fractions:

$$[3, 7, 15, 1, 292] = 3 + \frac{1}{7 + \frac{1}{15 + \frac{1}{1 + \frac{1}{292}}}}$$

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Infinite continued fractions:

$$[a_0, a_1, a_2, a_3, \ldots] = a_0 + rac{1}{a_1 + rac{1}{a_2 + rac{1}{a_3 + \cdots}}}$$

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$$\sqrt{389} = [19, 1, 2, 1, 1, 1, 1, 2, 1, 38, 1, 2, 1, 1, 1, 1, 2, 1, 38, \ldots]$$

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What about numbers of degree ≥ 3 ?

$$\sqrt[3]{2} = [1, 3, 1, 5, 1, 1, 4, 1, 1, 8, 1, 14, 1, 10, 2, 1, 4, 12, 2, 3, 2, 1]$$

 $3, 4, 1, 1, 2, 14, 3, 12, 1, 15, 3, 1, 4, 534, 1, 1, 5, 1, 1, \ldots]$
Theorem

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 $3, 4, 1, 1, 2, 14, 3, 12, 1, 15, 3, 1, 4, 534, 1, 1, 5, 1, 1, \ldots]$

Lang and Trotter: "no significant departure from random behaviour"

"[...] no continued fraction development of an algebraic number of higher degree than the second is known. It is not even known if such a development has bounded elements."



A. Khinchin. 1949.

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"Is there an algebraic number of degree higher than two whose simple continued fraction has unbounded partial quotients? Does every such number have unbounded partial quotients?"

R. K. Guy, 2004



Fact. The simple continued fraction expansion of $x \in \mathbb{R}$ is unbounded iff L(x) = 0.

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Theorem (Chonev, Ouaknine, W., 2015)

If the ZERO PROBLEM is decidable at order 9 then there is an algorithm that given a real algebraic number α computes $L(\alpha)$ to arbitrary precision. In particular, the set

 $\{\alpha \in \overline{Q} : \alpha \text{ has bounded partial quotients}\}$

would be recursively enumerable.

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Diophantine-approximation bounds play a key role in the proof—specifically Baker's theorem on linear forms in logarithms of algebraic numbers.



Consider the exponential polynomial

$$f(t) = 2 + \cos(t + \varphi_1) + \cos(\sqrt{2}t + \varphi_2) - e^{-t}$$

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Orbit $\{(t + \varphi_1, \sqrt{2}t + \varphi_2) \mod 2\pi : t \in \mathbb{R}_{\geq 0}\}$ is dense in $[0, 2\pi]^2$



Baker's Theorem:

$$\left|\left|\left(t+\varphi_1,\sqrt{2}t+\varphi_2\right)-(\pi,\pi)\right|\right|\geq \frac{1}{\operatorname{poly}(t)}$$

Conclusion and Perspectives

The Discrete Case

A linear recurrence sequence is a sequence $\langle u_0, u_1, u_2, \ldots \rangle$ of integers such that there exist constants a_1, \ldots, a_k , such that

$$u_{n+k} = a_1 u_{n+k-1} + a_2 u_{n+k-2} + \ldots + a_k u_n$$

for all $n \ge 0$.

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Theorem (Skolem 1934; Mahler 1935, 1956; Lech 1953)

The set of zeros of a linear recurrence sequence is semi-linear:

$$\{n: u_n = 0\} = F \cup A_1 \cup \ldots \cup A_\ell$$

where F is finite and each A_i is a full arithmetic progression.

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Theorem (Berstel and Mignotte 1976)

In Skolem-Mahler-Lech, the infinite part (arithmetic progressions A_1, \ldots, A_ℓ) is fully constructive.

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"... a mathematical embarrassment ... "

Richard Lipton

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- Even the bounded problem is hard.
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- Similar obstacles for the Infinite-Zeros Problem.
- Diophantine-approximation techniques unavoidable.